# DEPARTMENT OF <br> COMPUTER SCIENCE ENGINEERING 

III SEMESTER - R 2017

## CS8382 DIGITAL SYSTEMS LABORATORY

## LABORATORY MANUAL

## Name

Register No

Section

## VISION

is
committed to provide highly disciplined, conscientious and enterprising professionals conforming to global standards through value based quality education and training.

## MISSION

- To provide competent technical manpower capable of meeting requirements of the industry
- To contribute to the promotion of Academic Excellence in pursuit of Technical Education at different levels
- To train the students to sell his brawn and brain to the highest bidder but to never put a price tag on heart and soul


## DEPARTMENT OF COMPUTER SCIENCE ENGINEERING

## VISION

To provide candidates with knowledge and skill in the field of Electrical and Electronics Engineering and thereby produce extremely well trained employable, socially responsible and innovative Electrical and Electronics Engineers.

## MISSION

- To provide the students rigorous learning experience to produce creative solutions to society's needs.
- To produce electrical engineers of high calibre, conscious of the universal moral values adhering to professional ethical code.
- To provide highest quality learning environment for the students emphasizing fundamental concepts with strongly supported laboratory and prepare them to meet the global needs of the industry by continuous assessment and training.


## PROGRAM EDUCATIONAL OBJECTIVES (PEOs)

## 1. Fundamentals

To provide students with a solid foundation in mathematics, science and fundamentals of engineering enabling them to solve complex problems in order to develop real time applications.

## 2. Core Competence

To train the students to meet the needs of core industry with an attitude of learning new technologies.
3. Breadth

To provide relevant training and experience to bridge the gap between theory and practice which enable them to find solutions to problems in industry and research that contributes to the overall development of society.
4. Professionalism

To inculcate professional and effective communication skills to the students to make them lead a team and stand as a good decision maker to manage any constraint environment with good professional ethics at all strategies.

## 5. Lifelong Learning/Ethics

To practice ethical and professional responsibilities in the organization and society with commitment and lifelong learning needed for successful professional career.

## PROGRAM OUTCOMES (POs)

a. Graduates will demonstrate knowledge of mathematics, science and electrical engineering.
b. Graduates will be able to identify, formulate and solve electrical engineering problems.
c. Graduates will be able to design and conduct experiments, analyze and interpret data.
d. Graduates will be able to design a system, component or process as per needs and specifications.
e. Graduates will demonstrate to visualize and work on laboratory and multidisciplinary tasks.
f. Graduates will demonstrate skills to use modern engineering tools, software and equipment to analyze problems.
g. Graduates will demonstrate knowledge of professional and ethical responsibilities.
h. Graduates will be able to communicate effectively by both verbal and written form.
i. Graduates will show the understanding of impact of engineering solutions on the society and also will be aware of contemporary issues.
j. Graduates will develop confidence for self-education and ability for lifelong learning.
k. Graduate who can participate and succeed in competitive examinations.

## SYLLABUS

## Objectives:

## The student should be made to:

- Understand the various logic gates.
- Be familiar with various combinational circuits.
- Understand the various components used in the design of digital computers.
- Be exposed to sequential circuits
- Learn to use HDL

List of experiments:

1. Verification of Boolean Theorems using basic gates.
2. Design and implementation of combinational circuits using basic gates for arbitrary functions, code converters.
3. Design and implementation of combinational circuits using MSI devices:
a. 4 - bit binary adder / subtractor
b. Parity generator / checker
c. Magnitude Comparator
d. Application using multiplexers
4. Design and implementation of sequential circuits:
a. Shift -registers
b. Synchronous and asynchronous counters
5. Coding combinational / sequential circuits using HDL.
6. Design and implementation of a simple digital system (Mini Project).

## Course Outcomes:

- Use Boolean simplification techniques to design a combinational hardware circuit.
- Design and Implement combinational and sequential circuits.
- Analyze a given digital circuit - combinational and sequential.
- Design the different functional units in a digital computer system.
- Design and Implement a simple digital system.
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Content

| SI.No. | Name of the Experiment | Page No. |
| :---: | :--- | :--- |
| 1. | Verification of Boolean Theorems using Digital Logic Gates |  |
| 2. | Design and Implementation of Combinational Circuits using Basic Gates for Arbitrary <br> Functions, Code Converters |  |
| 3. | Implementation of half adder and full adder |  |
| 4. | Implementation of half subtractor and full subtractor |  |
| 5. | Design and Implementation of 4-Bit Binary Adder / Subtractor using Basic Gates and <br> MSI Devices |  |
| 6. | Design and Implementation of Parity Generator / Checker using Basic Gates and MSI <br> Devices |  |
| 7. | Design and Implementation of Magnitude Comparator. |  |
| 8. | Design and Implementation of Application using Multiplexers / Demultiplexers. |  |
| 9. | Design and Implementation of Shift Registers. |  |
| 10. | Design and Implementation of Synchronous and Asynchronous Counters. |  |
| 11. | Simulation of Combinational Circuits using Hardware Description Language (VHDL / Verilog <br> HDL Software Required). |  |
| 12. | Simulation of Sequential Circuits using HDL (VHDL / Verilog HDL Software Required). |  |

## Expt.No.1:

Aim:
To verify the truth table of basic digital IC's of AND, OR, NOT, NAND, NOR, EX-OR gates
Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :--- | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | NAND gate | IC 7400 | 1 |
| 6. | NOR gate | IC 7402 | 1 |
| 7. | EX-OR gate | IC 7486 | 1 |
| 8. | Connecting wires |  | As required |.

Theory:
Circuit that takes the logical decision and the process are called logic gates. Each gate has one or more input and only one output. OR, AND and NOT are basic gates. NAND, NOR and X-OR are known as universal gates. Basic gates form these gates.

## AND gate

The AND gate performs a logical multiplication commonly known as AND function. The output is high when both the inputs are high. The output is low level when any one of the inputs is low.

## OR gate

The OR gate performs a logical addition commonly known as OR function. The output is high when any one of the inputs is high. The output is low level when both the inputs are low.

## NOT gate

A NOT gate is the physical realization of the complementation operation. The NOT gate is called an inverter. The output is high when the input is low. The output is low when the input is high. NAND gate

The NAND gate is a contraction of AND-NOT. The output is high when both inputs are low and any one of the input is low. The output is low level when both inputs are high.
NOR gate
The NOR gate is a contraction of OR-NOT. The output is high when both inputs are low. The output is low when one or both inputs are high.

## EX-OR gate

An Ex-OR gate performs the following Boolean function, $A \in B=\left(A . B^{\prime}\right)+\left(A^{\prime} . B\right)$. It is similar to $O R$ gate but excludes the combination of both $A$ and $B$ being equal to one. The exclusive OR is a function that give an output signal ' 0 ' when the two input signals are equal either ' 0 ' or ' 1 '.

OR Gate:

OR GATE:


NOT Gate symbol:


TRUTH TABLE:

| $\mathbf{A}$ | $\bar{A}$ |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |

EXOR Gate symbol:


TRUTH TABLE :

| $\mathbf{A}$ | $\mathbf{B}$ | $\overline{\mathbf{A}} \mathbf{B}+\mathbf{A} \overline{\mathbf{B}}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

PIN Diagram:


PIN Diagram:


NAND Gate symbol:
$A-\underset{7400}{A-B}-Y=\overline{A \cdot B}$

TRUTH TABLE

| $A$ | $B$ | $\overline{A \cdot B}$ |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

PIN Diagram:


NOR Gate:

## Procedure:

1. Connections are given as per the circuit diagram.
2. For all the IC's $7^{\text {th }}$ pin is grounded and $14^{\text {th }}$ pin is given +5 supply.
3. Apply the inputs and verify the truth table for all gates.

## Result:

The truth tables of all the basic logic gates were verified.

## Outcome:

At the completion of an experiment student will able to verify the truth table of all basic gates

## Viva - Voce

1. List out the basic gate.
2. Mention the universal gate.
3. How many gates presented in IC 7408 ?
4. What is IC?
5. What are the applications of gates?
6. Write the truth table of AND gate.
7. Write the truth table of OR gate.
8. Write the truth table of NOT gate.
9. Write the truth table of NAND gate.
10. Write the truth table of NOR gate.
11. Write the truth table of EX- OR gate.
12. What are the classifications of IC?
13. What are types of linear integrated circuit?
14. What is meant by etching?
15. What are the advantages of IC?
16. Write the truth table of EX- NOR gate.

## Expt.No.2:

## VERIFICATION OF BOOLEAN THEOREMS USING LOGIC GATES

Aim: To verification of Boolean theorems using logic gates
Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | NAND gate | IC 7400 | 1 |
| 6. | NOR gate | IC 7402 | 3 |
| 7. | EX-OR gate | IC 7486 | 1 |
| 8. | Connecting wires |  | As required |

Theory:
BASIC Boolean Laws

1. Commutative Law

The binary operator OR, AND is said to be commutative if,

1. $A+B=B+A$
2. $A \cdot B=B \cdot A$
3. Associative Law

The binary operator OR, AND is said to be associative if,

1. $A+(B+C)=(A+B)+C$
2. $A \cdot(B \cdot C)=(A \cdot B) \cdot C$
3. Distributive Law

The binary operator OR, AND is said to be distributive if,

1. $A+(B . C)=(A+B) \cdot(A+C)$
2. $A .(B+C)=(A . B)+(A . C)$
3. Absorption Law
4. $A+A B=A$
5. $A+A B=\overline{A+} B$
6. Idempotent Law
7. $A+A=A$
8. $A \cdot A=A$
9. Complementary Law
10. $A+A^{\prime}=1$
11. $A . A^{\prime}=0$
12. De Morgan's Theorem
13. The complement of the sum is equal to the sum of the product of the individual complements.

$$
A+B=A \cdot B
$$

2. The complement of the product is equal to the sum of the individual complements.

$$
\overline{\mathrm{A} \cdot \mathrm{~B}}=\overline{\mathrm{A}}+\bar{B}
$$

Design

1. Absorption Law

$$
A+A B=A
$$


2. Involution (or) Double complement Law

$$
A=A
$$

3. Idempotent Law
4. $A+A=A$

5. $A \cdot A=A$

6. Demorgan's Law
$A+B=A \cdot B$

7. Distributive Law

$$
A+(B \cdot C)=(A+B) \cdot(A+C)
$$



## Procedure:

1. Obtain the required IC along with the Digital trainer kit.
2. Connect zero volts to $G N D$ pin and +5 volts to $V_{C C}$.
3. Apply the inputs to the respective input pins.
4. Verify the output with the truth table.

## Result:

Thus the above stated Boolean laws are verified.

## Outcome:

At the completion of an experiment student will able to know the basic laws with their truth table.
Viva - Voce

1. What is Demorgan's law?
2. What is associative law?
3. What is mean by compliment gate?
4. Explain the basic laws in digital electronics
5. What is double complement?

## Expt.No.3: HALF ADDER AND FULL ADDER

## Aim:

To design and verify the truth table of the Half Adder \& Full Adder circuits

## Apparatus required:

| S. No. | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | EX-OR gate | IC 7486 | 1 |
| 6. | Connecting wires | As required |  |

Theory:
The most basic arithmetic operation is the addition of two binary digits. There are four possible elementary operations, namely,

$$
\begin{aligned}
& 0+0=0 \\
& 0+1=1 \\
& 1+0=1 \\
& 1+1=0 \text { (with } 1 \text { as carry) }
\end{aligned}
$$

The first three operations produce a sum of whose length is one digit, but when the last operation is performed the sum is two digits. The higher significant bit of this result is called a carry and lower significant bit is called the sum.
Half adder:
A combinational circuit which performs the addition of two bits is called half adder. The input variables designate the augend and the addend bit, whereas the output variables produce the sum and carry bits.

Full adder:
A combinational circuit which performs the arithmetic sum of three input bits is called full adder. The three input bits include two significant bits and a previous carry bit. A full adder circuit can be implemented with two half adders and one OR gate.

From the truth table, the expression for sum and carry bits of the output can be obtained as,
$S U M=A^{\prime} B^{\prime} C+A^{\prime} B C^{\prime}+A B^{\prime} C^{\prime}+A B C$
$C A R R Y=A^{\prime} B C+A B^{\prime} C+A B C^{\prime}+A B C$

Truth table:

| SI.No. | Input |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
|  | A | B | S | C |
| 1. | 0 | 0 | 0 | 0 |
| 2. | 0 | 1 | 1 | 0 |
| 3. | 1 | 0 | 1 | 0 |
| 4. | 1 | 1 | 1 | 1 |

From the truth table the expression for sum and carry bits of the output can be obtained as, Sum,
$S=A \oplus B$
Carry, C = A . B

Circuit diagram:


Full adder
Truth table:

| SI.No. | Input |  |  | Output |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | A | B | C | S | C |
| 1. | 0 | 0 | 0 | 0 | 0 |
| 2. | 0 | 0 | 1 | 1 | 0 |
| 3. | 0 | 1 | 0 | 1 | 0 |
| 4. | 0 | 1 | 1 | 0 | 1 |
| 5. | 1 | 0 | 0 | 1 | 0 |
| 6. | 1 | 0 | 1 | 0 | 1 |
| 7. | 1 | 1 | 0 | 0 | 1 |
| 8. | 1 | 1 | 1 | 1 | 1 |

Using Karnaugh maps the reduced expression for the output bits can be obtained as,
Sum:


$$
S U M=A^{\prime} B^{\prime} C+A^{\prime} B C^{\prime}+A B^{\prime} C^{\prime}+A B C=A \oplus B \oplus C
$$

Carry:


$$
C A R R Y=A B+A C+B C
$$

Logic Diagram:


## Procedure:

1. Connections are given as per the circuit diagrams.
2. For all the IC 's $7^{\text {th }}$ pin is grounded and $14^{\text {th }}$ pin is given +5 V supply.
3. Apply the inputs and verify the truth table for the half adder and full adder circuits.

## Result:

The design of the half adder and full adder circuits was done and their truth tables were verified. Outcome:

At the completion of an experiment student will able to design the half adder circuit and the full adder circuit.
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Aim:
To design and verify the truth table of the half subtractor \& full subtractor circuits
Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | EX-OR gate | IC 7486 | 1 |
| 6. | Connecting wires | As required |  |

## Theory:

The subtraction of two binary digits has four possible operations. In all operations, each subtrahend bit is subtracted from the minuend bit. In case of the second operation the minuend bit is smaller than the subtrahend bit, hence 1 is borrowed.

Half subtractor:
A combinational circuit which performs the subtraction of two bits is called half subtractor. The input variables designate the minuend and the subtrahend bit, whereas the output variables produce the difference and borrow bits.

Full subtractor:
A combinational circuit which performs the subtraction of three input bits is called full subtractor. The three input bits include two significant bits and a previous borrow bit. A full subtractor circuit can be implemented with two half subtractors and one OR gate.

From the truth table the expression for difference and borrow bits of the output can be obtained as,

Difference, DIFF= $A^{\prime} B^{\prime} C+A^{\prime} B C^{\prime}+A B^{\prime} C^{\prime}+A B C$

Borrow, $B O R R=A^{\prime} B C+A B^{\prime} C+A B C^{\prime}+A B C$

Truth table:

| SI.No. | Input |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
|  | A | B | Difference | Borrow |
| 1. | 0 | 0 | 0 | 0 |
| 2. | 0 | 1 | 1 | 1 |
| 3. | 1 | 0 | 1 | 0 |
| 4. | 1 | 1 | 0 | 0 |

From the truth table the expression for difference and borrow bits of the output can be obtained as,
Difference, DIFF =A $\oplus B$
Borrow, BORR = A'. B

Logic diagram:

2. Full subtractor

Truth table:

|  | Input |  |  |  | Output |
| :---: | :---: | :---: | :---: | :---: | :---: |
| SI.No. | A | B | C | Difference | Borrow |
| 1. | 0 | 0 | 0 | 0 | 0 |
| 2. | 0 | 0 | 1 | 1 | 1 |
| 3. | 0 | 1 | 0 | 1 | 1 |
| 4. | 0 | 1 | 1 | 0 | 1 |
| 5. | 1 | 0 | 0 | 1 | 0 |
| 6. | 1 | 0 | 1 | 0 | 0 |
| 7. | 1 | 1 | 0 | 0 | 0 |
| 8. | 1 | 1 | 1 | 1 | 1 |

Using Karnaugh maps the reduced expression for the output bits can be obtained as,

Difference


Borrow


$$
\text { Borrow }=A^{\prime} B+A^{\prime} C+B C
$$

Circuit diagram:


Expt.No.5:

Aim:
To design and implement 4-bit adder and subtractor using IC 7483
Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | IC | IC 7483 | 1 |
| 3. | NOT gate | IC 7404 | 1 |
| 4. | EX-OR gate | IC 7486 | 1 |
| 5. | Connecting wires |  | As required |

## Theory:

## 4 BIT Binary adder:

A binary adder is a digital circuit that produces the arithmetic sum of two binary numbers. It can be constructed with full adders connected in cascade, with the output carry from each full adder connected to the input carry of next full adder in chain. The augends bits of ' $A$ ' and the addend bits of ' $B$ ' are designated by subscript numbers from right to left, with subscript 0 denoting the least significant bits. The carries are connected in chain through full adder. The input carry to the adder is $\mathrm{C}_{0}$ and it ripples through the full adder to the output carry C4.

4 BIT Binary subtractor:
The circuit for subtracting A-B consists of an adder with inverters, placed between each data input ' B ' and the corresponding input of full adder. The input carry C 0 must be equal to 1 when performing subtraction.

## 4 BIT Binary adder / subtractor:

The addition and subtraction operation can be combined into one circuit with one common binary adder. The mode input $M$ controls the operation. When $M=0$, the circuit is adder circuit. When $M=1$, it becomes
subtractor.

PIN Diagram for IC 7483:

| 1 |  | A4 |  | B4 |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | - | 53 | I | 54 |  |
| 3 | - | A3 | C | C4 |  |
| 4 | - | B3 | 7 | C1 |  |
| 5 | - | VCC | 4 | GND |  |
| 6 | - | 52 | 8 | B1 |  |
| 7 | - | B2 | 3 | A1 |  |
| 8 | - | A2 |  | S1 |  |

Logic Diagram: 4-Bit Binary Diagram:


Logic diagram: 4-Bit Binary Subtractor:


4-Bit Binary Adder /Subtractor:

$\mathrm{M}=\mathbf{0}$ [ADDITION] $\mathrm{M}=1$ [SUBTRACTION]

Truth table:

| Input Data A |  |  |  | Input Data B |  |  |  | Addition |  |  |  |  | Subtraction |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A4 | A3 | A2 | A1 | B4 | B3 | B2 | B1 | C | S4 | S3 | S2 | S1 | B | D4 | D3 | D2 | D1 |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 1 |
| 1 | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 0 | 1 | 1 | 1 | 0 | 1 | 1 | 0 | 1 |

## Procedure:

1. Connections are given as per the circuit diagrams.
2. Logical inputs were given as per circuit diagram.
3. Apply the inputs and verify the truth table for the 4-bit adder and subtractor.

## Result:

The design of the 4-bit Binary adder and I subtractor circuit was done and its truth table was verified.

Outcome:
At the completion of an experiment student will able to design 4-bit binary adder and subtractor

## Viva - Voce

1. What is expression for difference and borrow?
2. Write the truth table for half adder.
3. Write the truth table for full adder.
4. Write the truth table for half subtrator.
5. Write the truth table for full subtrator.
6. Draw the logic diagram of full subtrator.
7. What is adder?
8. List out the application of adders.
9. Draw the full adder using two half adder circuits.
10. What is combinational circuit?
11. What is different between combinational and sequential circuit?
12. What are the gates involved for binary adder?
13. List the properties of Ex-Nor gate?
14. What is expression for sum and carry?

## Expt.No.6:

Aim:
To design, construct and study the performance of 2 bit magnitude comparator
Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | Magnitude comparator | IC 7485 | 2 |
| 6. | EX-OR gate | IC 7486 | 1 |
| 7. | Connecting wires |  | As required |

## Theory:

The comparison of two numbers is an operator that determines one number is greater than, less than (or) equal to the other number. A magnitude comparator is a combinational circuit that compares two numbers $A$ and $B$ and determines their relative magnitude. The outcome of the comparator is specified by three binary variables that indicate whether $A>B, A=B$ (or) $A<B$.

Truth table:

| Inputs |  |  |  |  | Outputs |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{A}_{1}$ | $\mathbf{A}_{0}$ | $\mathbf{B}_{1}$ | $\mathbf{B}_{0}$ | $\mathbf{A}>\mathbf{B}$ | $\mathbf{A}=\mathbf{B}$ | $\mathbf{A}<\mathbf{B}$ |  |
| 0 | 0 | 0 | 0 | 0 | 1 | 0 |  |
| 0 | 0 | 0 | 1 | 0 | 0 | 1 |  |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 |  |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 |  |
| 0 | 1 | 0 | 0 | 1 | 0 | 0 |  |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 |  |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 |  |
| 0 | 1 | 1 | 1 | 0 | 0 | 1 |  |
| 1 | 0 | 0 | 0 | 1 | 0 | 0 |  |
| 1 | 0 | 0 | 1 | 1 | 0 | 0 |  |
| 1 | 0 | 1 | 0 | 0 | 1 | 0 |  |
| 1 | 0 | 1 | 1 | 0 | 0 | 1 |  |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 |  |
| 1 | 1 | 0 | 1 | 1 | 0 | 0 |  |
| 1 | 1 | 1 | 0 | 1 | 0 | 0 |  |
| 1 | 1 | 1 | 1 | 0 | 1 | 0 |  |

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K-map

$\mathrm{A}>\mathrm{B}=\mathrm{A} 0 \overline{\mathrm{~B} 0} \mathrm{~B} 1+\mathrm{A} 1 \overline{\mathrm{~B} 1}+\mathrm{A} 1 \mathrm{~A} 0 \overline{\mathrm{~B} 0}$



Logic Diagram:


Pin Diagram for IC 7485:

| B3 | -1 |  | 16 |  |
| :---: | :---: | :---: | :---: | :---: |
| 1(A<B) | - 2 | I | 15 |  |
| 1(A=B) | $-3$ | C | 14 |  |
| 1(A>B) | - 4 | 7 | 13 |  |
| A>B | - 5 | 4 | 12 |  |
| $A=B$ | -6 | 8 | 11 |  |
| A<B | -7 | 5 | 10 |  |
| GND | -8 |  | 9 |  |

## 8 Bit Magnitude Comparator:



Truth table:

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{A}>\mathbf{B}$ | $\mathbf{A}=\mathbf{B}$ | $\mathbf{A}<\mathbf{B}$ |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{0 0 0 0}$ | $\mathbf{0 0 0 0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{0 0 0 0}$ | $\mathbf{0 0 0 0}$ |  |  |  |
| $\mathbf{0 0 0 1}$ | $\mathbf{0 0 0 0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{0 0 0 1}$ | $\mathbf{0 0 0 0}$ |  |  |  |
| $\mathbf{0 0 0 0}$ | $\mathbf{0 0 0 1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{0 0 0 0}$ | $\mathbf{0 0 0 1}$ |  |  |  |

Procedure:

1. Connections are given as per circuit diagram.
2. Logical inputs are given as per circuit diagram.
3. Observe the output and verify the truth table.

## Result:

Thus the 2-bit and 8-bit magnitude comparator was designed and verified using the logic gates.

## Outcome:

At the completion of an experiment student will able to design the 2-bit and 8-bit magnitude comparator using logic gates.

1. What is magnitude comparator?
2. What is most significant bit?
3. Explain operation of AND gate.
4. Explain truth table of a comparator.
5. Explain magnitude comparator7485 IC.
6. What is 8 -bit input Magnitude Comparator?
7. What is IC?
8. Explain the $k$-map simplification of $A>B$.
9. Explain the $k$-map simplification of $A=B$.
10. Explain the $k$-map simplification of $A<B$
11. Draw the logic diagram of 1-bit magnitude comparator.
12. What is the truth table of 1 -bit magnitude comparator?
13. What is the use of magnitude comparator?
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## Aim:

To design, construct and study the performance of 4-bit different code converters
(i) Binary to gray code converter
(ii) Gray to binary code converter
(iii) BCD to excess-3 code converter
(iv) Excess-3 to BCD code converter

## Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :--- | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | Magnitude comparator | IC 7485 | 2 |
| 6. | EX-OR gate | IC 7486 | 1 |
| 7. | Connecting wires |  | As required |

## Theory:

The availability of large variety of codes for the same discrete elements of information results in the use of different codes by different systems. A conversion circuit must be inserted between the two systems if each uses different codes for same information. Thus, code converter is a circuit that makes the two systems compatible even though each uses different binary code. The bit combination assigned to binary code to gray code. Since each code uses four bits to represent a decimal digit. There are four inputs and four outputs. Gray code is a non-weighted code. The input variable are designated as B3, B2, B1, B0 and the output variables are designated as $\mathrm{C} 3, \mathrm{C} 2, \mathrm{C} 1, \mathrm{Co}$. from the truth table, combinational circuit is designed. The Boolean functions are obtained from K-Map for each output variable. A code converter is a circuit that makes the two systems compatible even though each uses a different binary code. To convert from binary code to Excess-3 code, the input lines must supply the bit combination of elements as specified by code and the output lines generate the corresponding bit combination of code. Each one of the four maps represents one of the four outputs of the circuit as a function of the four input variables. A twolevel logic diagram may be obtained directly from the Boolean expressions derived by the maps. These are various other possibilities for a logic diagram that implements this circuit. Now the OR gate whose output is C+D has been used to implement partially each of three outputs.
(i) Binary to gray code converter

Logic Diagram:


K map for G3:


K map for G2:

$\mathrm{G} 2=\mathrm{B} 3 \oplus \mathrm{~B} 2$


K map for G 0 :


Truth table:

|  |  |  |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |

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Logic Diagram:


K map for B 3 :


K map for B2:

$\mathrm{B} 2=\mathrm{G} 3 \oplus \mathrm{G} 2$

K map for B1:

$\mathrm{B} 1=\mathrm{G} 3 \oplus \mathrm{G} 2 \oplus \mathrm{G} 1$

K map for B 0 :


Truth table:

| G3 | G2 | G1 | G0 | B3 | B2 | B1 | B0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 0 | 1 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 1 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 | 1 | 1 | 1 | 1 |

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(iii) BCD to excess-3 code converter

Logic Diagram:


K map for E3:


K map for E 2 :


K map for E1:


K map for E :


$$
\mathrm{E} 0=\overline{\mathrm{B} 0}
$$

Truth table:

| B3 | B2 | B1 | B0 | G3 | G2 | G1 | G0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ | $\mathbf{x}$ |
|  |  |  |  |  |  |  |  |

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BCD code converter
Logic Diagram:


K map for A :

| 00 | 00 | 01 | 11 | 10 |
| :---: | :---: | :---: | :---: | :---: |
|  | X | X | 0 | X |
| 01 | 0 | 0 | 0 | 0 |
| 11 | 1 | X | X | X |
| 10 | 0 | 0 | 1 | 0 |

K map for B:


K map for C :


K map for D:


Truth table:

| B3 | B2 | B1 | B0 | G3 | G2 | G1 | G0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |  |  |  |
| 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 |
| 1 | 0 | 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 1 | 1 | 0 | 0 | 0 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |

## Procedure:

1. Connections were given as per circuit diagram.
2. Logical inputs were given as per truth table
3. Observe the logical output and verify with the truth tables.

## Result:

Thus the code converters were designed and verified using the corresponding truth table.

## Outcome:

At the completion of an experiment student will able to design the binary to gray converter.

## Viva - Voce

1. What is binary code?
2. What is gray code?
3. What are the advantages of gray code?
4. What is unit distance code?
5. What is sequential code?
6. How to convert binary to gray code?
7. How to convert gray to binary code?
8. What is reflective code?
9. What are the advantages of $E X-3$ code?
10. Which code is used to arithmetic operation in digital circuits?
11. Explain the operation of $E X-O R$.
12. What is K - Map?
13. Draw the truth table of EX- OR gate.
14. What is SOP?
15. What is POS?
16. What is minterm?

## Expt.No.8: PARITY GENERATORS AND CHECKERS

Aim:
To implement the odd and even parity checkers using the logic gates and also to generate the odd parity and even parity numbers using the generators

## Apparatus required:

| Sl. No | Component | Type | Quantity |
| :---: | :--- | :---: | :---: |
| 1 | Trainer Kit | - | 1 |
| 2 | EX-OR | IC7486 | 1 |
| 3 | NOT gate | IC 7404 | 1 |
| 4 | Connecting wires | - | Required |

## Theory:

Parity checking is used for error detection in data transmission.
Odd parity checkers:
It counts the number of 1 's in the given input and produces a 1 in the output when the number of 1 's is odd.

Even parity checker:
It counts the number of 1 's in the given input and produces a 1 in the output when the number of 1 's is even.

Odd parity generators:
It generates an odd parity number. The odd parity checker circuit is used with the inverted output and also the input bits. So when the input is a 4-bit number then the output of the generator circuit will have 5 bits which is an odd parity number.

Even parity generator:
It generates an even parity number. The even parity checker circuit is used with the inverted output and also the input bits. So when the input is a 4-bit number then the output of the generator circuit will have 5 bits which is an even parity number.

| Input |  |  |  |  | Checker output |  | Generator output |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | B | C | D | odd | even | odd | even |  |
| 0 | 0 | 0 | 1 | 1 | 0 | 00010 | 00011 |  |
| 0 | 0 | 1 | 0 | 1 | 0 | 00100 | 00101 |  |
| 0 | 0 | 1 | 1 | 0 | 1 | 00111 | 00110 |  |
| 0 | 1 | 0 | 0 | 1 | 0 | 01000 | 01001 |  |
| 0 | 1 | 0 | 1 | 0 | 1 | 01011 | 01010 |  |
| 0 | 1 | 1 | 0 | 0 | 1 | 01101 | 01100 |  |
| 0 | 1 | 1 | 1 | 1 | 0 | 01110 | 01111 |  |
| 1 | 0 | 0 | 0 | 1 | 0 | 10000 | 10001 |  |
| 1 | 0 | 0 | 1 | 0 | 1 | 10011 | 10010 |  |
| 1 | 0 | 1 | 0 | 0 | 1 | 10101 | 10100 |  |
| 1 | 0 | 1 | 1 | 1 | 0 | 10110 | 10111 |  |
| 1 | 1 | 0 | 0 | 0 | 1 | 11001 | 11000 |  |
| 1 | 1 | 0 | 1 | 1 | 0 | 11010 | 11011 |  |
| 1 | 1 | 1 | 0 | 1 | 0 | 11100 | 11101 |  |
| 1 | 1 | 1 | 1 | 0 | 1 | 11111 | 11110 |  |

## Procedure:

1. The circuit is implemented using logic gates.
2. The inputs are given as per the truth table.
3. The corresponding outputs are noted.
4. The theoretical and practical values were verified.

## Result:

The odd and even parity checkers are implemented using the logic gates and the odd parity and even parity numbers are generated using the corresponding generators.

## Outcome:

At the completion of an experiment student will able to verify the odd and even parity checker using logic gates.

1. What is parity bit?
2. Why parity bit is added to message?
3. What is parity checker?
4. What is odd parity?
5. What is even parity?
6. What are the gates involved for parity generator?
7. List the procedures to convert gray code into binary.
8. Why weighted code is called as reflective codes?
9. What is a sequential code?
10. What is error deducting code?
11. What is ASCII code?
12. What is hamming code?
13. List the binary weighted code.
14. List the binary non weighted code.
15. Write the hamming code equation
16. List the procedures to convert binary code into gray
17. What are the applications of gray code?
18. What are the applications of Excess- 3 code?

## Expt.No.9:

Aim:
To design and verify the truth table of a 4X1 multiplexer \& 1X4 demultiplexer
Apparatus required:

| SI. No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | OR gate | IC 7432 | 1 |
| 3. | NOT gate | IC 7404 | 1 |
| 4. | AND gate ( three input ) | IC 7411 | 1 |
| 5. | Connecting wires |  | As required |

## Theory:

Multiplexing means transmitting a large number of information units over a smaller number of channels or lines. A digital multiplexer is a combinational circuit that selects binary information from one of many input lines and directs it to a single output line. The selection of particular input line is controlled by a set of selection lines. Normally, there are $2^{n}$ input lines and $n$ selection lines whose bit combinations determines which input is selected. A multiplexer is called a data selector, since it selects one of many inputs and steers the binary information to the output line. A Strobe is also provided to allow the designer to disable all output data until a specified time. Then, by allowing the STROBE to go low, the proper lead can be selected. This feature is very useful where data might be changing the same time DATA SELECT leads change. It is a very useful Medium Scale Integration (MSI) function and has a multitude of applications. It is used for connecting two or more sources to a single destination among the computer units and it is useful for constructing a common bus system. A decoder with an enable input can function as a demultiplexer. A Demultiplexer is a circuit that receives information on a single line and transmits this information on one of $2^{n}$ possible output lines. The selection of specific output line is controlled by the bit values of $n$ selection lines. The decoder and demultiplexer operations are obtained from the same circuit; a decoder with an enable input is referred to as a decoder / de-multiplexer. The Strobe lead can be used to active or deactive the entire IC, allowing time for the address lines to change the information is fed to the output. Demultiplexers are useful anytime information from one source must be fed several places.


CIRCUIT DIAGRAM:


## Result:

The design of the $4 \times 1$ Multiplexer and $1 \times 4$ Demultiplexer circuits was done and their truth tables were verified.

Outcome:
At the completion of an experiment student will able to design the multiplexer and the demultiplexer
Viva - Voce

1. What is multiplexer?
2. What is demultiplexer?
3. What are the advantages of multiplexer?
4. What are the advantages of demultiplexer?
5. What is select signal?
6. How to choose select signal in multiplexer?
7. How to choose select signal in demultiplexer?
8. Write the formula used in select signal.
9. What is the difference between the multiplexer and demultiplexer?
10. What is the application of multiplexer?
11. What is the application of demultiplexer?
12. Draw the truth table of multiplexer.
13. Draw the truth table of demultiplexer.
14. How many select signals are needed in $8 \star 1$ multiplexer?
15. How many select signals are needed in $8 \star 1$ demultiplexer?
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Aim:

To design and implement the various shift register
Apparatus required:

| SI. No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | D flip flop | IC 7474 | 2 |
| 2. | OR gate | IC 7432 | 1 |
| 3. | IC Trainer kit |  | 1 |
| 5. | Connecting wires | As required |  |.

Theory:

A register is capable of shifting its binary information in one or both directions is known as shift register. The logical configuration of shift register consist of a D-Flip flop cascaded with output of one flip flop connected to input of next flip flop. All flip flops receive common clock pulses which causes the shift in the output of the flip flop. The simplest possible shift register is one that uses only flip flop. The output of a given flip flop is connected to the input of next flip flop of the register. Each clock pulse shifts the content of register one bit position to right.

## PIN Diagram:



## Logic Diagram:

SERIAL IN SERIAL OUT


| CLK | Serial in | Serial out |
| :---: | :---: | :---: |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| 2 | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{3}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{4}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{5}$ | $\mathbf{X}$ | $\mathbf{0}$ |
| $\mathbf{6}$ | $\mathbf{X}$ | $\mathbf{0}$ |
| 7 | $\mathbf{X}$ | $\mathbf{1}$ |

Logic Diagram:
Serial in parallel out:


Truth Table:

| CLK | DATA | OUTPUT |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $Q_{A}$ | $Q_{B}$ | $Q_{C}$ | $Q_{D}$ |
| 1 | 1 | 1 | 0 | 0 | 0 |
| 2 | 0 | 0 | 1 | 0 | 0 |
| 3 | 0 | 0 | 0 | 1 | 1 |
| 4 | 1 | 1 | 0 | 0 | 1 |

Parallel in Serial Out:

## Truth Table:

| CLK | Q3 | Q2 | Q1 | Q0 | $\mathbf{O} / \mathbf{P}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{2}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ |
| $\mathbf{3}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |

Parallel in Parallel Out:

PARALLEL IN PARALLEL OUT:

CLR

## Truth Table:

| $\mathbf{C L K}^{\|c\|}$ | DATA INPUT $^{\prime}$ |  |  |  | OUTPUT |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathbf{D}_{\mathbf{A}}$ | $\mathbf{D}_{\mathbf{B}}$ | $\mathbf{D}_{\mathbf{C}}$ | $\mathbf{D}_{\mathbf{D}}$ | $\mathbf{Q}_{\mathbf{A}}$ | $\mathbf{Q}_{\mathbf{B}}$ | $\mathbf{Q}_{\mathbf{C}}$ | $\mathbf{Q}_{\mathbf{D}}$ |
|  | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{2}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |

## Procedure:

1. Connections are given as per circuit diagram
2. Logical inputs are given as per circuit diagram.
3. Observe the output and verify the truth table.

Result:
Thus the implementation of shift registers using flip flops was completed successfully.
Outcome:
At the completion of an experiment student will able to design the various types of shift register.

## Expt．No．12：SYNCHRONOUS UPIDOWN COUNTER

Aim：
To design and implement 3 bit synchronous up／down counter
Apparatus required：

| S．No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :---: | :---: |
| 1. | JK Flip Flop | IC 7474 | 2 |
| 2. | OR gate | IC 7432 | 1 |
| 3. | NOT gate | IC 7404 | 1 |
| 4. | AND gate（ three input ） | IC 7411 | 1 |
| 5 | XOR gate | IC 7486 | 1 |
| 6 | IC Trainer Kit |  | 1 |
| 7. | Connecting wires |  | As required |

Theory：
A counter is a register capable of counting number of clock pulse arriving at its clock input． Counter represents the number of clock pulses arrived．An up／down counter is one that is capable of progressing in increasing order or decreasing order through a certain sequence．An up／down counter is also called bidirectional counter．Usually up／down operation of the counter is controlled by up／down signal．When this signal is high counter goes through up sequence and when up／down signal is low counter follows reverse sequence．

K map：

Characteristic Table:

| $\mathbf{Q}$ | $\mathbf{Q}_{\mathbf{t} \mathbf{1}}$ | $\mathbf{J}$ | $\mathbf{K}$ |
| :---: | :---: | :---: | :---: |
| $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{X}$ |
| $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{X}$ |
| $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{X}$ | $\mathbf{1}$ |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{X}$ | $\mathbf{0}$ |

Logic Diagram:

## Truth Table:

| Input Up/Down | Present <br> State |  |  | $\begin{gathered} \text { Next State } \\ \mathbf{Q}_{\mathrm{A}+1} \mathrm{Q}_{\mathrm{B}+1} \mathrm{Q}_{\mathrm{C}+1} \end{gathered}$ |  |  | A |  | B |  | C |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | $\mathrm{J}_{\mathrm{A}}$ | $\mathbf{K}_{\text {A }}$ |  | $\mathrm{K}_{\mathrm{B}}$ | $\mathrm{J}_{\mathbf{C}}$ | $\mathbf{K}_{\text {C }}$ |
| 0 | 0 | 0 | 0 |  |  |  | 1 | 1 | 1 | 1 | X | 1 | X | 1 | X |
| 0 | 1 | 1 | 1 | 1 | 1 | 0 | X | 0 | X | 0 | X | 1 |
| 0 | 1 | 1 | 0 | 1 | 0 | 1 | X | 0 | X | 1 | 1 | X |
| 0 | 1 | 0 | 1 | 1 | 0 | 0 | X | 0 | 0 | X | X | 1 |
| 0 | 1 | 0 | 0 | 0 | 1 | 1 | X | 1 | 1 | X | 1 | X |
| 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 | X | X | 0 | X | 1 |
| 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 | X | X | 1 | 1 | X |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | X | 0 | X | X | 1 |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | X | 0 | X | 1 | X |
| 1 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | X | 1 | X | X | 1 |
| 1 | 0 | 1 | 0 | 0 | 1 | 1 | 0 | X | X | 0 | 1 | X |
| 1 | 0 | 1 | 1 | 1 | 0 | 0 | 1 | X | X | 1 | X | 1 |
| 1 | 1 | 0 | 0 | 1 | 0 | 1 | X | 0 | 0 | X | 1 | X |
| 1 | 1 | 0 | 1 | 1 | 1 | 0 | X | 0 | 1 | X | X | 1 |
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | X | 0 | X | 0 | 1 | X |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | X | 1 | X | 1 | X | 1 |

## Procedure:

1. Connections are given as per circuit diagram.
2. Logical inputs are given as per circuit diagram.
3. Observe the output and verify the truth table.

Result:
Thus the 3-bit synchronous up/down counters was implemented successfully.
Outcome:
At the completion of an experiment student will able to design the synchronous up/down counter.

## Expt.No.12:

## SIMULATION OF COMBINATIONAL CIRCUITS USING HDL

Aim:
To write a verilog code for half adder, full adder and multiplexer
Tools Required:
Xilinx 9.2

Program:

Simulation wave for half adder

## Procedure:

1. Write and draw the Digital logic system.
2. Write the Verilog code for above system.
3. Enter the Verilog code in Xilinx software.
4. Check the syntax and simulate the above Verilog code (using ModelSim or Xilinx) and verify the output waveform as obtained.
5. Implement the above code in Spartan III using FPGA kit.

## Result:

Thus the verilog code for half adder, full adder and multiplexer were simulated and verified successfully.

## Outcome:

At the completion of an experiment student will able to be known the verilog code for half adder, full adder and multiplexer.

## Expt.No.13:

## SIMULATION OF SEQUENTIAL CIRCUITS USING HDL

Aim:
To write a verilog code for RS, D, JK flip flop and up counter
Tools required:
Xilinx 9.2
Program:
RS Flip Flop:

```
module sr1(clk, s, r, q);
    input clk;
    input s;
    input r;
    output reg q;
    always@(posedge clk or s or r)
    begin
    if(clk==1)
    if(s==1 && r==0)
    q<=1;
    else if(s==0 && r==0)
    q<=q;
    else
    q<=0;
    end
endmodule
```



D Flip Flop:
module dffi (d, clk, rst, q);
input $d$;
input clk;
input rst;
output q;
reg q;
always@ (posedge clk)
if(rst)
q<=0;
else
q<=d;
endmodule


JK Flip Flop:

```
module jkflip(clk, j, k, q);
    input clk;
    input j;
    input k;
    output reg q;
    always@ (posedge clk or j or k)
    begin
    if(clk==1)
    if(j==1 && k==0)
    q<=1;
    else if(j==0 && k==0)
    q<=q;
    else if(j==0 && k==1)
    q<=0;
    else
    q<=~q;
    end
endmodule
```



Up Counter:

```
module upcount (clk, N, a);
    input clk;
    input [3:0] N;
    output reg [3:0] a;
    initial a=4'b0000;
    always @(negedge clk)
    a=(a==N) 24'b0000: a+1'b1;
endmodule
```



## Result:

Thus the verilog code for RS,D,JK Filp Flop and up counter were simulated and verified successfully.

## Outcome:

At the completion of an experiment student will able to be known the verilog code for RS, D, JK Filp Flop and up counter .

## ADDITIONAL EXPERIMENTS

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## Expt.No.8: ENCODER AND DECODER

Aim:

To study the operation of encoder and decoder circuits using logic gates

Apparatus required:

| S. No | Name of the Apparatus |  | Range | Quantity |
| :---: | :--- | :--- | :--- | :---: |
| 1. | Digital IC trainer |  |  | 1 |
| 2. | NOT Gate | IC 7404 |  | 1 |
| 3. | OR Gate | IC 7432 |  | 1 |
| 4. | AND Gate | IC7408 |  | 1 |
| 5. | Bread Board | IC7404 |  | 1 |
| 6. | NOT Gate |  |  | 1 |
| 8. | Connecting wires and probes | As required |  |  |

## Theory:

## Decoder

In digital electronics, a decoder can take the form of a multiple-input, multiple-output logic circuit that converts coded inputs into coded outputs, where the input and output codes are different e.g. n-to-2n, binary-coded decimal decoders. Decoding is necessary in applications such as data multiplexing, 7 segment display and memory address decoding.

The example decoder circuit would be an AND gate because the output of an AND gate is "High" (1) only when all its inputs are "High." Such output is called as "active High output". If instead of AND gate, the NAND gate is connected the output will be "Low" (0) only when all its inputs are "High". Such output is called as "active low output".

A slightly more complex decoder would be the $n$-to-2n type binary decoders. These types of decoders are combinational circuits that convert binary information from ' n ' coded inputs to a maximum of 2 n unique outputs. In case the 'n' bit coded information has unused bit combinations, the decoder may have less than $2 n$ outputs. 2-to-4 decoder, 3-to-8 decoder or 4-to-16 decoder are other examples.

The input to a decoder is parallel binary number and it is used to detect the presence of a particular binary number at the input. The output indicates presence or absence of specific number at the decoder input. An encoder is a device, circuit, transducer, software program, algorithm or person that converts information from one format or code to another. The purpose of encoder is standardization, speed, secrecy, security, or saving space by shrinking size. Encoders are combinational logic circuits and they are exactly opposite of decoders. They accept one or more inputs and generate a multibit output code. Encoders perform exactly reverse operation than decoder. An encoder has M input and Noutput lines. Out
of $M$ input lines only one is activated at a time and produces equivalent code on output $N$ lines. If a device output code has fewer bits than the input code has, the device is usually called an encoder

|  | Inputs |  | Outputs |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SI.No. | A | B | Y 3 | Y 2 | Y1 | Y0 |
| 1. | 0 | 0 | 0 | 0 | 0 | 1 |
| 2. | 0 | 1 | 0 | 0 | 1 | 0 |
| 3. | 1 | 0 | 0 | 1 | 0 | 0 |
| 4. | 1 | 1 | 1 | 0 | 0 | 0 |


| SI.No. | Inputs |  |  |  |  |  |  |  |  |  |  |  |  |  | Outputs |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | D7 | D6 | D 5 | D 4 | D 3 | D 2 | D 1 | D 0 | A | B | C |  |  |  |  |  |  |
| 1. | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |  |  |  |  |  |  |
| 2. | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 |  |  |  |  |  |  |
| 3. | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |  |  |  |  |  |  |
| 4. | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |  |  |  |  |  |  |
| 5. | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |  |  |  |  |  |  |
| 6. | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 |  |  |  |  |  |  |
| 7. | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |  |  |  |  |  |  |
| 8. | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 |  |  |  |  |  |  |

Procedure:

1. Make the circuit connections as shown in the figure.
2. Check the corresponding truth table.

## Result:

The design of the Encoder and Decoder circuit was done and the input and output were obtained

## Outcome:

At the completion of an experiment student will able to design the encoder circuit and the decoder circuit

1. What is Encoder?
2. What is decoder?
3. List the application of encoder.
4. List the application of decoder.
5. Draw the truth table of encoder.
6. Draw the truth table of decoder.
7. What are logic gates used encoder?
8. What are logic gates used encoder?
9. What is the difference between decoder with demultiplexer?
10. What is the difference between encoder with multiplexer?
11. How to choose the select signal in encoder?
12. How to choose the select signal in decoder?
13. Draw the logic diagram of encoder.
14. Draw the logic diagram of encoder.
15. What is the difference between encoder with decoder?

## Expt.No.2: IMPLEMENTATION OF BOOLEAN FUNCTIONS

Aim:
To design the logic circuit and verify the truth table of the given Boolean expression,

$$
F(A, B, C, D)=\Sigma(0,1,2,5,8,9,10)
$$

Apparatus required:

| SI.No | Name of the Apparatus | Range | Quantity |
| :---: | :--- | :--- | :---: |
| 1. | Digital IC trainer kit |  | 1 |
| 2. | AND gate | IC 7408 | 1 |
| 3. | OR gate | IC 7432 | 1 |
| 4. | NOT gate | IC 7404 | 1 |
| 5. | NAND gate | IC 7400 | 1 |
| 6. | NOR gate | IC 7402 | 1 |
| 7. | EX-OR gate | IC 7486 | 1 |
| 8. | Connecting wires |  | As required |

Circuit diagram:


Design:
Given, $F(A, B, C, D)=\Sigma(0,1,2,5,8,9,10)$
Truth table:

| SI. No. | INPUT |  |  |  | OUTPUT |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | A | B | C | D | $\mathrm{F}=\mathrm{D}^{\prime} \mathrm{B}^{\prime}+\mathrm{C}^{\prime}\left(\mathrm{B}^{\prime}+\mathrm{A}^{\prime} \mathrm{D}\right)$ |
| 1. | 0 | 0 | 0 | 0 | 1 |
| 2. | 0 | 0 | 0 | 1 | 1 |
| 3. | 0 | 0 | 1 | 0 | 1 |
| 4. | 0 | 0 | 1 | 1 | 0 |
| 5. | 0 | 1 | 0 | 0 | 0 |
| 6. | 0 | 1 | 0 | 1 | 1 |
| 7. | 0 | 1 | 1 | 0 | 0 |
| 8. | 0 | 1 | 1 | 1 | 0 |
| 9. | 1 | 0 | 0 | 0 | 1 |
| 10. | 1 | 0 | 0 | 1 | 1 |
| 11. | 1 | 0 | 1 | 0 | 1 |
| 12. | 1 | 0 | 1 | 1 | 0 |
| 13. | 1 | 1 | 0 | 0 | 0 |
| 14. | 1 | 1 | 0 | 1 | 0 |
| 15. | 1 | 1 | 1 | 0 | 0 |
| 16. | 1 | 1 | 1 | 1 | 0 |

The output function F has four input variables hence a four variable Karnaugh Map is used to obtain a simplified expression for the output as shown,


From the K-Map,
$F=B^{\prime} C^{\prime}+D^{\prime} B^{\prime}+A^{\prime} C^{\prime} D$
Since we are using only two input logic gates the above expression can be re-written as,
$F=C^{\prime}\left(B^{\prime}+A^{\prime} D\right)+D^{\prime} B^{\prime}$
Now the logic circuit for the above equation can be drawn.

## Procedure:

1. Connections are given as per the circuit diagram.
2. For all the IC 's $7^{\text {th }}$ pin is grounded and $14^{\text {th }}$ pin is given +5 V supply.
3. Apply the inputs and verify the truth table for the given Boolean expression.

## Result:

The truth table of the given Boolean expression was verified.

## Outcome:

At the completion of an experiment student will able to design the Boolean expression.

