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Sl. No. of Ques. Paper : 1332

Unique Paper Code

: 2341502

Name of Paper Name of Course

: Theory of Computation : B.Tech. Computer Science

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Semester

Duration:

: 3 hours

Maximum Marks

: 75

(Write your Roll No. on the top immediately on receipt of this question paper.)

Question No. 1 is of 35 marks and all its parts are compulsory. Attempt four questions from Q. Nos 2 to 7.

## PART A

Note: For all the questions, consider the alphabet {a,b} unless otherwise specified.

(a) Is (S\*)\* = (S\*)\* for all sets S? Explain with an example.

(b) Consider the language PALINDROME and a string y over the given alphabet. Prove that if the string y3 is in PALINDROME, then so is the string y.

(3)

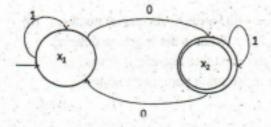
(c) Give a regular expression for the language of all words that do not end in (3) a double letter.

(d) Show that (ab)\*a and a(ba)\* define the same language. Give the set of strings representing the two languages. Give the first five strings generated in the lexicographic manner.

- (e) Using pumping lemma for regular languages, show that the language. (3)  $L=\{a^nba^n \mid n \ge 0\}$  is not regular.
  - (4)

(f) Given two Finite Automata(FA): FA1 and FA2 find the machine for the intersection of the languages represented by these FA's.

FAI

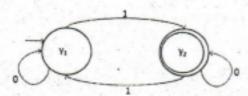


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- (g) Create a Push Down Automata(PDA) for the language L = { a<sup>n</sup>S, where S (4) starts with b and length (S) = n }.
- (h) Find a Context Free Grammar(CFG) for the language defined by the regular expression a\*b\*.
  (2)
- Show that the following CFG is ambiguous by finding a word with two distinct syntax trees:

S-AA

A-AAA|a|bA |Ab

(j) Convert the following CFG into CNF: (3)

S-aXX

X-aS | bS | a

(k) Explain the working of the following Turing Machine(TM) (2)

 $>R \xrightarrow{a\neq U} R \xrightarrow{b\neq U} R_U a R_U b$ 

U represents the blank symbol.

(l) Describe the Universal Turing Machine.

(4)

## PART B

2(a) Let language L<sub>1</sub>= EQUAL, the language with words having equal number of a's and b's and L<sub>2</sub>= {a<sup>n</sup>b<sup>m</sup>a<sup>n</sup> | m,n=1,2,...}. What is the language defined by the intersection of L1 and L2? Is it a context free language? If yes, construct a PDA for the language, else prove using pumping lemma for CFLs.

- 3(a) Prove that regular languages are closed under complementation, i.e., if a (3) language L is regular then L'(complement of L) is also regular.
- 3 (b) For the following pair of regular languages find an FA that defines the (4) difference, L<sub>1</sub>-L<sub>2</sub>:

L1=(a+b) c

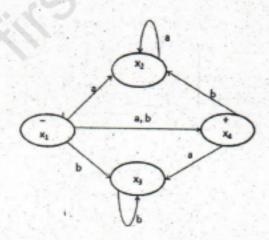
L2=b(a+b) c

 $\Sigma = \{a,b,c\}$ 

- 3 (c) Build an FA that accepts the language of all strings of a's and b's such that (3) the next to last letter is an a.
- 4(a) Consider the homomorphism h from the alphabet {0,1,2} to {a,b}.defined (4) by:

h(0)=ab, h(1)=b, h(2)=aa

- i) What is h(0210)?
- ii) What is h(2201)?
- iii) If L is the language 1 02, what is h(L)?
- 4(b) Give a PDA for the language with words of type  $a^xb^ya^zb^wx$ , y, z, w = 1, 2,3 (6) ... y>x and z>w and x+z=y+w.
- 5(a) Convert the following NFA to DFA. (5)



5 (b) Write regular expression and construct a DFA for the following language (5) of all words that have an even number of substrings ab in them.

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6(a) Consider the following CFG in Chomsky Normal Form (CNF) (6) S-PO Q-QS|b Generate the derivation trees for the word i)abab ii)ababab Consider S as the self embedded non terminal, trace the division of each word w into uvxyz and uvvxyyz, where |u| + |z| >= 0, |v| + |y| > 0 and |x| > 0. 6(b) Which of the following could be configurations of a Turing Machine? Justify your answer. i.

- (q, ≥aUaU, U, Ua)
  - ii. (q, abc,b, abc)
  - iii. (p, ≥ abc, a, e) (h, ▶, e, e)

(represents the left end symbol)

- 7 (a) Give a Turing Machine which computes the function f(w) =ww. (5)
- 7 (b) The language H = { "M""w" ; Turing machine M halts on input w} (5)describes the halting problem. Prove that H is not recursive, i.e., the Halting problem is undecidable.